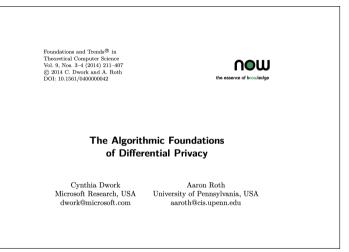
# Trustworthy Machine Learning Differential Privacy 1

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POSTECH

### **Contents from**



• and contents partially from Gautam Kamath at University of Waterloo and Roger Grosse at University of Toronto.

## Why Privacy Guarantees in Learning?

• Not anonymized dataset for learning - privacy leak

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- Not anonymized dataset for learning privacy leak
- Anonymized dataset for learning looks okay but possible to leak private information

# Why Privacy Guarantees in Learning?

**Anonymized Dataset** 

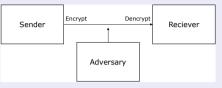
Name	Age	Gender	Zip Code	Smoker	Diagnosis
*	60-70	Male	191**	Y	Heart disease
*	60-70	Female	191**	N	Arthritis
*	60-70	Male	191**	Y	Lung cancer
*	60-70	Female	191**	N	Crohn's disease
*	60-70	Male	191**	Y	Lung cancer
*	50-60	Female	191**	N	HIV
*	50-60	Male	191**	Y	Lyme disease
*	50-60	Male	191**	Y	Seasonal allergies
*	50-60	Female	191**	N	Ulcerative colitis

Figure: An example from Kearns & Roth, The Ethical Algorithm

- anonymized dataset looks okay but still privacy leak
  - If we know Rebecca is 55 years old and in this database, then we know she has 1 of 2 diseases.

# Why Not Use Cryptosystems?

### Set-up for Encryption



• Entities in encryption: Sender, Receiver, and Adversary

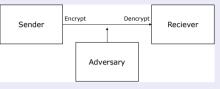
### Set-up for Private Learning



- e.g., a learning algorithm (=curator) releases a model for some "benefits"
- But, the model should not reveal private information.

# Why Not Use Cryptosystems?

### Set-up for Encryption



• Entities in encryption: Sender, Receiver, and Adversary

### Set-up for Private Learning



- e.g., a learning algorithm (=curator) releases a model for some "benefits"
- But, the model should not reveal private information.
- Note that homomorphic encryption could be alternatives but slow (yet)

## **Goal for Privacy In Learning**

### Goal

Learning nothing about an individual while learning useful information about a population.

• How to achieve this goal? Add noise!

## **Goal for Privacy In Learning**

### Goal

Learning nothing about an individual while learning useful information about a population.

- How to achieve this goal? Add noise!
- In the learning context,
  - > Here, an algorithm can transform a dataset into another dataset
  - Sender: an algorithm that releases a model
  - Reciever: A model user

An Example

### Goal of A Survey

Estimate a statistic on illegal behaviors of participants, where

- Curator: a participant
- Analyst: a researcher

An Example

### Goal of A Survey

Estimate a statistic on illegal behaviors of participants, where

- Curator: a participant
- Analyst: a researcher
- Each participant follows the following survey process:
  - Flip a coin
  - If "tails", then respond truthfully.
  - () If "heads", then flip a second coin and respond "Yes" if "heads" and "No" if "tails".

### **General Description**

Randomized Response

$$Y_i = \begin{cases} X_i & \text{with probability } \frac{1}{2} + \gamma \\ 1 - X_i & \text{with probability } \frac{1}{2} - \gamma, \end{cases}$$

- $X_i \in \{0,1\}$ : the truthful response
- $Y_i \in \{0,1\}$ : a randomized response
- $\gamma = 0$ : a uniformly random strategy
  - 🗸 private
  - ✗ not informative
- $\gamma = 1/2$ : an honest strategy
  - 🗡 no privacy
  - $\checkmark$  informative
- $\gamma = 1/4$ : the previous example.
  - $\checkmark\,$  private  $\rightarrow\,$  no learning on an individual response
  - $\checkmark$  informative  $\rightarrow$  learning on a population statistic

#### How Informative?

Randomized Response

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- Observe the following expectation over "unfair coin flips":

$$\mathbb{E}_{\mathcal{Q}}\{Y_i\} = X_i\left(\frac{1}{2} + \gamma\right) + (1 - X_i)\left(\frac{1}{2} - \gamma\right) = 2\gamma X_i + \frac{1}{2} - \gamma \implies X_i = \mathbb{E}_{\mathcal{Q}}\left\{\frac{1}{2\gamma}\left(Y_i - \frac{1}{2} + \gamma\right)\right\}$$

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• Consider the following estimator:

$$\hat{p} \coloneqq \frac{1}{n} \sum_{i=1}^{n} \left( \frac{1}{2\gamma} \left( Y_i - \frac{1}{2} + \gamma \right) \right)$$

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• The randomized response looks "working"! How can we prove that this "algorithm" does not leak privacy? 9/25

# A Goodness Metric in Differential Privacy (DP)

### Definition

A randomized algorithm  $\mathcal{M}$  is  $(\varepsilon, \delta)$ -differentially private if for any  $\mathcal{S} \in \text{Range}(\mathcal{M})$  and for any two "neighboring" datasets  $\mathcal{D}_1$  and  $\mathcal{D}_2$ ,

$$\mathbb{P} \left\{ \mathcal{M}(\mathcal{D}_1) \in \mathcal{S} \right\} \le \exp(\varepsilon) \mathbb{P} \left\{ \mathcal{M}(\mathcal{D}_2) \in \mathcal{S} \right\} + \delta,$$

where the probability is taken over the randomness of  $\mathcal{M}$ .

• Consider the following special case (*i.e.*,  $\delta = 0$  and  $\varepsilon \to 0$ ):

$$1 \approx \frac{1}{\exp(\varepsilon)} \le \frac{\mathbb{P}\left\{\mathcal{M}(\mathcal{D}_1) \in \mathcal{S}\right\}}{\mathbb{P}\left\{\mathcal{M}(\mathcal{D}_2) \in \mathcal{S}\right\}} \le \exp(\varepsilon) \approx 1$$

- After applying differentially private  $\mathcal{M}$ , *i.e.*,  $\mathcal{S} = \mathcal{M}(\mathcal{D}_1)$ , an attacker cannot tell whether  $\mathcal{S}$  is from  $\mathcal{D}_1$  or  $\mathcal{D}_2$  so cannot extract information from the difference between  $\mathcal{D}_1$  and  $\mathcal{D}_2$ .
- e.g.,  $\mathcal{D}_1 = \{X_1 = 0, X_2 = 1\}$ ,  $\mathcal{D}_2 = \{X_1 = 0\}$ ,  $\mathcal{S} = \{1\}$ ,  $\mathcal{M} =$  "contain 1?"

## Randomized Response is DP

### Theorem

The randomized response is  $(\ln 3, 0)$ -differentially private.

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### Proof sketch.

- $\mathcal{M}$ : a randomized response
  - $\mathcal{M}(X_1,\ldots,X_n) = (Y_1,\ldots,Y_n)$
- Let  $\gamma = \frac{1}{4}$
- Consider any realization  $\mathcal{S} \in \{0,1\}^n$  of  $(Y_1, \ldots, Y_n)$ .
- Consider  $X \coloneqq (X_1, \dots, X_n)$  and  $X' \coloneqq (X'_1, \dots, X'_n)$  which differ only in coordinate j.
- Then, we have

$$\frac{\mathbb{P}\{\mathcal{M}(X)=\mathcal{S}\}}{\mathbb{P}\{\mathcal{M}(X')=\mathcal{S}\}} = \frac{\prod_{i=1}^{n} \mathbb{P}\{\mathcal{M}(X_i)=\mathcal{S}_i\}}{\prod_{i=1}^{n} \mathbb{P}\{\mathcal{M}(X'_i)=\mathcal{S}_i\}} = \frac{\mathbb{P}\{\mathcal{M}(X_j)=\mathcal{S}_j\}}{\mathbb{P}\{\mathcal{M}(X'_j)=\mathcal{S}_j\}} = \frac{\mathbb{P}\{Y_j=\mathcal{S}_j\}}{\mathbb{P}\{Y'_j=\mathcal{S}_j\}} \le \frac{1/2+\gamma}{1/2-\gamma} = e^{\ln 3}$$

• Note that  $S_j$  is fixed but the left-hand side of the inequality maximizes if  $S_j = 1$ .

## Laplace Mechanism

### Definition

Given any function  $f: \mathbb{N}^{|\mathcal{X}|} \to \mathbb{R}^k$ , the Laplace mechanism is defined as:

$$\mathcal{M}_L(x, f, \varepsilon) \coloneqq f(x) + (Y_1, \dots, Y_k),$$

where  $Y_i$  are i.i.d. random variables drawn from Lap  $\left(f(x)_i \mid \frac{\Delta f}{\varepsilon}\right)$ .

• Lap
$$(x|b) = Lap(b) = \frac{1}{2b} \exp\left(-\frac{|x|}{b}\right)$$

• The  $\ell_1\text{-sensitivity of a function }f:\mathbb{N}^{|\mathcal{X}|}\to\mathbb{R}^k$  is

$$\Delta f \coloneqq \max_{x,y \in \mathbb{N}^{|\mathcal{X}|}, \|x-y\|_1 = 1} \|f(x) - f(y)\|_1.$$

• e.g., x is a dataset and f is a post-processor.

### Laplace Mechanism is DP

### Theorem

The Laplace mechanism preserves  $(\varepsilon, 0)$ -differential privacy.

## Laplace Mechanism is DP

**Proof Sketch** 

- Let  $x \in \mathbb{N}^{|\mathcal{X}|}$  and  $y \in \mathbb{N}^{|\mathcal{X}|}$  be such that  $\|x y\|_1 \leq 1$
- $p_x$ : the PDF of  $\mathcal{M}_L(x, f, \varepsilon)$ , *i.e.*,  $p_x(z) \coloneqq \mathbb{P}\{\mathcal{M}_L(x, f, \varepsilon) = z\}$
- $p_y$ : the PDF of  $\mathcal{M}_L(y, f, \varepsilon)$
- For any  $z \in \mathbb{R}^k$ , we have

$$\frac{p_x(z)}{p_y(z)} = \prod_{i=1}^k \left( \exp\left(-\frac{\varepsilon |f(x)_i - z_i|}{\Delta f}\right) \middle/ \exp\left(-\frac{\varepsilon |f(y)_i - z_i|}{\Delta f}\right) \right)$$
$$= \prod_{i=1}^k \exp\left(\frac{\varepsilon (|f(y)_i - z_i| - |f(x)_i - z_i|)}{\Delta f}\right)$$
$$\leq \prod_{i=1}^k \exp\left(\frac{\varepsilon |f(x)_i - f(y)_i|}{\Delta f}\right)$$
$$= \exp\left(\frac{\varepsilon ||f(x) - f(y)||_1}{\Delta f}\right)$$
$$\leq \exp\left(\varepsilon\right).$$

•  $\frac{p_x(z)}{p_y(z)} \ge \exp(-\varepsilon)$  follows by symmetry.

### Gaussian Mechanism is DP

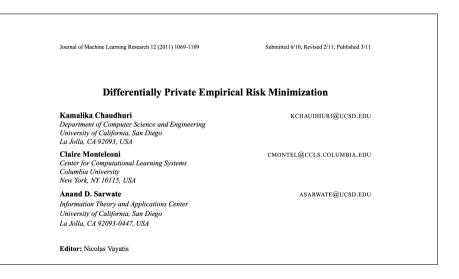
### Definition

Let  $f: \mathbb{N}^{|\mathcal{X}|} \to \mathbb{R}^d$  be an arbitrary *d*-dimensional function, and define its  $\ell_2$  sensitivity to be  $\Delta_2 f = \max_{\text{adjacent}x,y} \|f(x) - f(y)\|_2$ . The Gaussian Mechanism with parameter  $\sigma$  adds noise scaled to  $\mathcal{N}(0, \sigma^2)$  to each of the *d* components of the output.

### Theorem

Let  $\varepsilon \in (0,1)$  be arbitrary. For  $c^2 > 2\ln\left(\frac{1.25}{\delta}\right)$ , the Gaussan Mechanism with parameter  $\sigma \geq \frac{c\Delta_2 f}{\varepsilon}$  is  $(\varepsilon, \delta)$ -differentially private.

### How Can It be Connected to Learning?



### Empirical Risk Minimization (ERM) Setup

- $\mathcal{X}$ : an example space
  - Assume that  $\|x\|_2 \leq 1$  for  $x \in \mathcal{X}$
- $\bullet \ \mathcal{Y}:$  a label space
- $\mathcal{D}\coloneqq \{(x_i,y_i)\}_{i=1}^n \subseteq \mathcal{X} \times \mathcal{Y}$ : a training set
- $f: \mathcal{X} \to \mathcal{Y}$ : a predictor
- $\ell:\mathcal{Y}\times\mathcal{Y}\rightarrow\mathbb{R}:$  a loss function
- Regularized empirical risk minimization (ERM):

$$J(f, \mathcal{D}) \coloneqq \frac{1}{n} \sum_{i=1}^{n} \ell(f(x_i), y_i) + \Lambda N(f),$$

where N(f) is a regularizer.

## Assumptions

### Definition

A function f(x) over  $x \in \mathbb{R}^d$  is said to be strictly *convex* if for all  $\alpha \in (0, 1)$ , x, and  $y \neq x$ ,

$$f(\alpha x + (1 - \alpha)y) < \alpha f(x) + (1 - \alpha)f(y).$$

It is said to be  $\lambda\text{-strongly convex}$  if for all  $\alpha\in(0,1),$   $x\text{, and }y(\neq x)\text{,}$ 

$$f(\alpha x + (1 - \alpha)y) \le \alpha f(x) + (1 - \alpha)f(y) - \frac{1}{2}\lambda\alpha(1 - \alpha)\|x - y\|_{2}^{2}.$$

- A strictly convex function has a unique minimum.
- $(strongly convex) \Longrightarrow (strictly convex)$
- $\bullet$  The regularizer  $N(\cdot)$  and loss  $\ell(\cdot, \cdot)$  are differentiable.
  - ▶ No ℓ<sub>1</sub>-norm regularizer
  - No hinge loss
- Note that these assumptions are for handy analyses (and could be relaxed).

## **Privacy Model**

**Goal:** Learn a classifier which preserves the privacy of individual entities of a training set  $\mathcal{D}$ .

### Definition ( $\varepsilon$ -differential privacy)

An algorithm  $\mathcal{A}$  provides  $\varepsilon$ -differential privacy if for any two data sets  $\mathcal{D}$  and  $\mathcal{D}'$  that differ in a single entry and for any  $\mathcal{S}$ 

$$e^{-\varepsilon} \leq \frac{\mathbb{P}\{\mathcal{A}(\mathcal{D}) \in \mathcal{S}\}}{\mathbb{P}\{\mathcal{A}(\mathcal{D}') \in \mathcal{S}\}} \leq e^{\varepsilon}.$$

- $\mathcal{A}(\mathcal{D})$ : a randomized algorithm that returns a classifier from a training set  $\mathcal{D}$ .
- $\mathcal{D}'$  and  $\mathcal{D}$  have n-1 samples  $(x_i, y_i)$  in common; the different sample contains private values.

## Is ERM differentially private?

 $\bullet$  Given  ${\mathcal D}$  and  ${\mathcal D}',$  let

$$f_{\mathcal{D}}^* = \arg\min_f J(f, \mathcal{D}) \text{ and } f_{\mathcal{D}'}^* = \arg\min_f J(f, \mathcal{D}')$$

- Letting  $S := \{f_D^*\}$ ,  $\mathbb{P}\{f_D^* \in S\} = 1 \neq \mathbb{P}\{f_{D'}^* \in S\} = 0$ 
  - Note that our ERM is deterministic.
- Thus, ERM is not differentially private!

## **Algorithm 1: Output Perturbation**

**Output** Perturbation

$$f_{\mathsf{priv}} = \arg\min_{f} J(f, \mathcal{D}) + \mathbf{b}$$

 ${\ensuremath{\, \bullet \,}}$  b is random noise with density

$$v(\mathbf{b}) \propto e^{-\beta \|\mathbf{b}\|}$$

with  $\beta = \frac{n\Lambda\varepsilon}{2}$ .

• This algorithm is randomized.

## **Algorithm 2: Objective Perturbation**

**Objective Perturbation** 

$$f_{\mathsf{priv}} = rg\min_{f} J(f, \mathcal{D}) + rac{1}{n} \mathbf{b}^T f$$

 $\bullet~{\bf b}$  is random noise with density

$$v(\mathbf{b}) \propto e^{-\beta \|\mathbf{b}\|}$$

with 
$$\beta = \frac{\varepsilon - \log\left(1 + \frac{2c}{n\Lambda} + \frac{c^2}{n^2\Lambda^2}\right)}{2}$$
 (assuming  $\varepsilon$  is chosen to be  $\beta > 0$ ).

• This algorithm is randomized.

## **Privacy Guarantee**

#### Theorem

If  $N(\cdot)$  is differentiable and 1-strongly convex, and  $\ell$  is convex and differentiable with  $|\ell'(z)| \leq 1$  for all z, then Algorithm 1 provides  $\varepsilon$ -differential privacy.

### Theorem

If  $N(\cdot)$  is doubly differentiable and 1-strongly convex, and  $\ell$  is convex and doubly differentiable with  $|\ell'(z)| \leq 1$  and  $|\ell''(z)| \leq c$  for all z, then Algorithm 2 provides  $\varepsilon$ -differential privacy.

- Algorithm 2 requires stronger assumptions.
- What's the benefit of Algorithm 2?

### **Correctness Guarantee**

#### Lemma

Suppose  $N(\cdot)$  is doubly differentiable with  $\|\nabla N(f)\|_2 \leq \eta$  for all  $f, \ell$  is differentiable and has continuous c-Lipschitz derivatives. Given  $\mathcal{D}$ , let  $f^* := \arg \max_f J(\mathcal{D}, f)$  let  $f_{priv}$  be the output of Algorithm 1. Then, we have

$$\mathbb{P}_{\mathbf{b}}\left\{J(f_{\textit{priv}}, \mathcal{D}) - J(f^*, \mathcal{D}) \leq \frac{2d^2\left(\frac{c}{\Lambda} + \eta\right)\log^2\frac{d}{\delta}}{\Lambda n^2\varepsilon^2}\right\} \geq 1 - \delta$$

#### Lemma

Suppose  $N(\cdot)$  is 1-strongly convex and globally differentiable, and  $\ell$  is convex and differentiable with  $|\ell'(z)| \leq 1$  for all z. Given  $\mathcal{D}$ , let  $f^* := \arg \max_f J(\mathcal{D}, f)$  and let  $f_{priv}$  be the output of Algorithm 2. Then, we have

$$\mathbb{P}_{\mathbf{b}}\left\{J(f_{\text{priv}}, \mathcal{D}) - J(f^*, \mathcal{D}) \leq \frac{4d^2 \log^2 \frac{d}{\delta}}{\Lambda n^2 \varepsilon^2}\right\} \geq 1 - \delta.$$

- If  $\frac{c}{\Lambda} + \eta > 2$ , Algorithm 2 is better.
- Intuition: if perturbations are considered in learning, the algorithm finds a better classifier.

### Conclusion

- Differential privacy in learning:
  - $\blacktriangleright$  Hide "local" information  $\rightarrow$  satisfying the privacy guarantee
  - $\blacktriangleright$  Learn "global" information  $\rightarrow$  satisfying the correctness guarantee
- Two goals are conflicting each other and balancing two is critical.